

Equilibrium Selection in global games with strategic complementarities

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Dicussion by Michal Szkup

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- We refer to the above game $G(v)$ as a global game.

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- Focus on the games with strategic complementarities;

Assumptions

① (A1) Strategic Complementarities:

If $a_i \geq a'_i$ and $a_{-i} \geq a'_{-i}$ then $\forall \theta, \Delta u_i(a_i, a'_i, a_{-i}, \theta) \geq \Delta u_i(a_i, a'_i, a'_{-i}, \theta)$

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② (A2) Dominance regions

$\exists \underline{\theta}$ such that $\forall \theta < \underline{\theta}$ and $\forall a'_i \neq 0, \Delta u(0, a'_i, a_{-i}, \theta) > 0$

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- ③ (A3) State monotonicity:

$\exists K_0$ such that $\forall a_i \geq a'_i$ and $\theta, \theta' \in [\underline{\theta}, \bar{\theta}]$, $\theta \geq \theta'$ we have
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- ④ (A4) Payoff continuity

$u_i(a_i, a_{-i}, \theta)$ is continuous in all arguments

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- A mixed strategy is a probability distribution over pure strategies

Limit Uniqueness

Theorem (1)

\exists an increasing strategy profile s^* such that if, for each $v > 0$, s^v is a strategy profile that survives iterative deletion of strategies in $G(v)$, then

$$\lim_{v \rightarrow 0} s_i^v(x_i) = s_i^*(x_i)$$

for almost all x_i .

This theorem states that as noise vanishes the iterative deletion of dominated strategies selects a unique Bayesian Nash Equilibrium of the game

Intuition for Theorem 1 (Simple Case)

- 2 players with $A_i = [0, 1]$ and uniform prior

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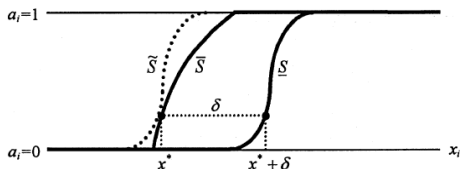
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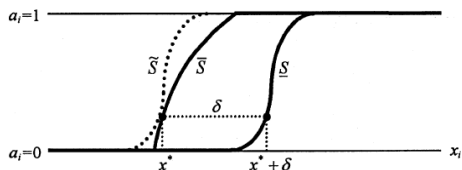
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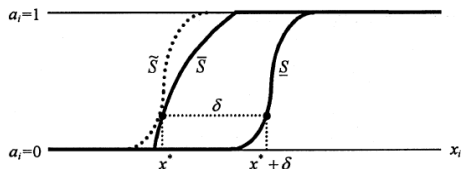
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- $a^* \geq \overline{S}(x^*) = \underline{S}(x^* + \delta)$ and if $\delta > 0$ $\underline{S}(x^* + \delta) > a^*$
- $\implies \delta = 0$ and so $\overline{S} = \underline{S} = \tilde{S}$

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- More precisely, the set of rationalizable strategies of the simplified game and original game converges.
- This result holds for any prior and any finite number of players or continuum of players.

A partial characterization of equilibrium

We can characterize the surviving equilibria of the game when noise is small

Definition

$Q(\varepsilon, \nu)$ is the set of parameters θ for which no Nash Equilibrium action profile a of the complete information game with payoffs $(u_i(\cdot, \theta'))_{i=1}^I$ for some $\theta' \in [\theta - \varepsilon, \theta + \varepsilon]$, such that for every strategy s^ν surviving iterative deletion of dominated strategies in $G(\nu)$, $\forall i \quad |s_i^\nu(\theta) - a_i| < \varepsilon$

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- Define $Q(\varepsilon, \nu)$ be a set of θ for which surviving strategy profiles in $G(\nu)$ do not prescribe all players to play close to some pure strategy Nash equilibrium of a complete information game.

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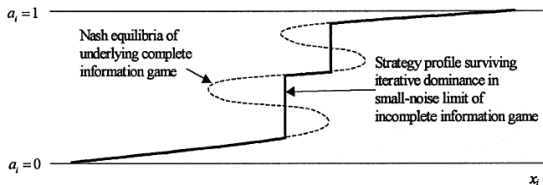
Theorem (2)

In $G(v)$ in the limit as $v \rightarrow 0$ for almost all payoff parameter θ , players play arbitrary close to some pure strategy Nash Equilibrium of the complete information game with payoffs $u(\cdot, \theta')$ that is arbitrary close to θ , i.e. $\forall \epsilon > 0$ there is \bar{v} s.t. $\forall v < \bar{v}$, $|\theta' - \theta| < \epsilon$

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Further results

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There exists a two-person, four-action game satisfying (A1) – (A5) in which for different noise structure different equilibria are selected in the limit as the signal errors vanish.

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Theorem (4)

If the complete information game at some payoff parameter θ is quasiconcave in a_i and has local potential maximizer (LP-maximizer) a^ then $s^*(\theta) = a^*$ regardless of the noise structure*

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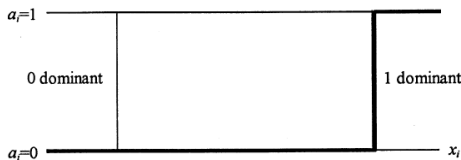
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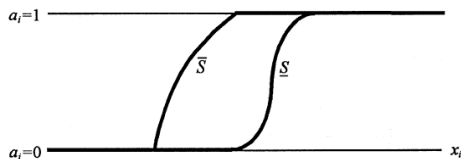
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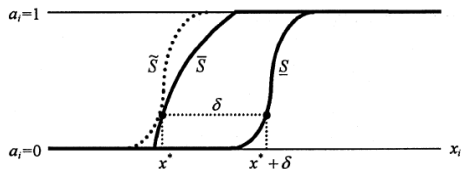
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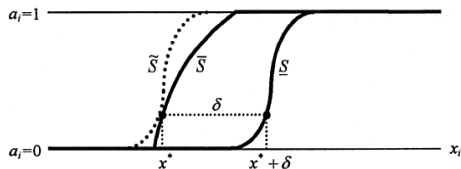
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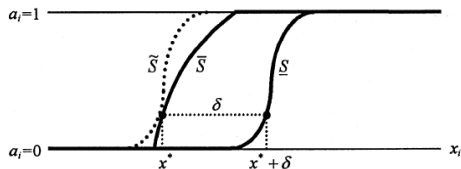
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- However, his estimate of θ is strictly higher at $x^* + \delta$
- by (A3) it follows that $\forall \delta > 0 \underline{S}(x^* + \delta) > a^* \implies \delta = 0$

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- However, his estimate of θ is strictly higher at $x^* + \delta$
- by (A3) it follows that $\forall \delta > 0 \underline{S}(x^* + \delta) > a^* \implies \delta = 0$
- This last implication follows from the fact that we also have $a^* \geq \underline{S}(x^* + \delta)$

Appendix: Final Step (2)

- Consider two cases:
 - ▶ player i observes x^* and expects his opponent to play \tilde{S}
 - ▶ player i observes $x^* + \delta$ and expects his opponent to play \underline{S}
- In both cases he faces the same distribution of action of his opponent.
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- by (A3) it follows that $\forall \delta > 0 \underline{S}(x^* + \delta) > a^* \implies \delta = 0$
- This last implication follows from the fact that we also have $a^* \geq \underline{S}(x^* + \delta)$
- That would lead to contradiction unless $\tilde{S}(x) = \underline{S}(x)$

Appendix (Potential Games)

Definition (Potential Function)

Let N be a finite set of players, Y_i be the set of i 's strategies and $u_i : Y \rightarrow R$ be i 's payoff function. Then a function $P : Y \rightarrow R$ is called a potential function if for every i and every Y_{-i} we have

$$u(x, y_{-i}) - u(x', y_{-i}) = P(x, y_{-i}) - P(x', y_{-i})$$

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- A potential function is a common payoff function such that a change in player i 's payoff from switching action to another is always the same as the change in the potential function.

Appendix (Intuition for Theorem 4)

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- a game is symmetric and the set of action is finite

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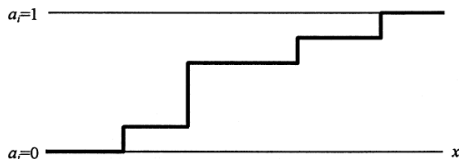
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 - ▶ this iterative procedure converges to equilibrium
 - ▶ but a_x^* is payoff maximizing profile, so as noise vanishes the best responses must be except for small neighborhood of discontinuity points
 - ▶ hence as noise vanishes, players must play potential maximizing action